PUBLIC KEY CRYPTO WITH APPLICATIONS

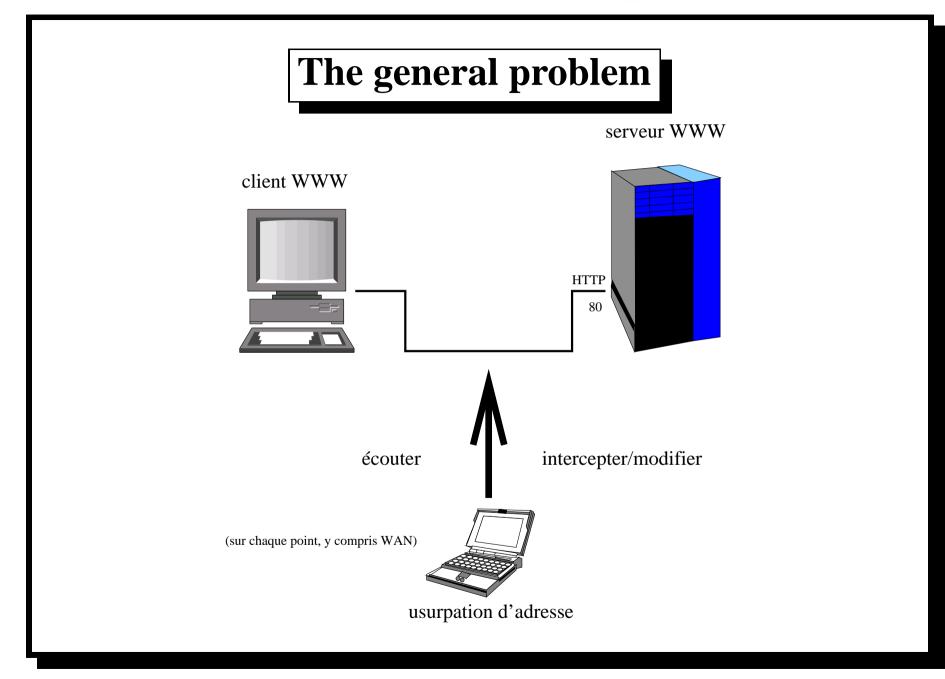
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Programme

- The general problem
- The available tools
- The concepts
- The problem of Trust
- Other attacks
- PKI
- Exchange your keys !
- Questions

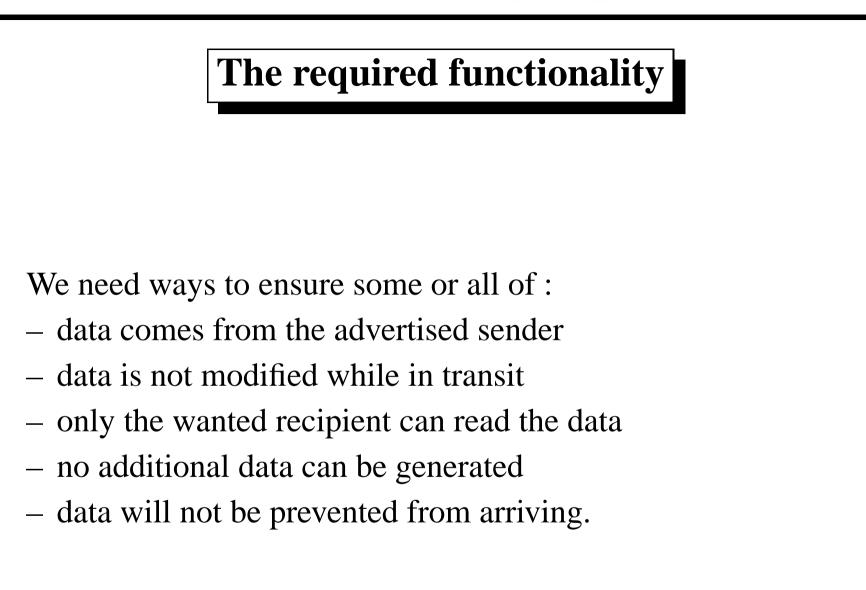


The problem

- The communication channel isn't safe (in this example : Internet and the HTTP protocol are not safe).
- Attackers can prevent data from arriving, modify them, generate new data and read them.

Targets

- "New" ways of using Internet
- electronic payment
- personal/medical data
- cracking
- social engineering

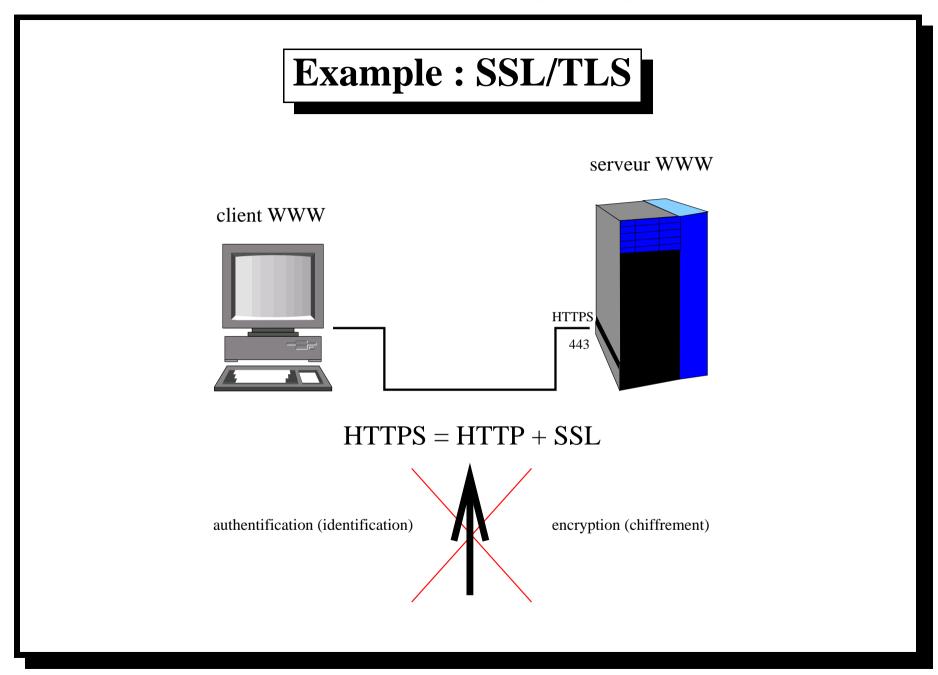


What users want

- ensure received data are reliable/true
- ensure personal/medical data stay confidential
- pay securely over the Internet (Yellownet, Internet banking)
- other future and sometimes irrealist applications : e-voting, e-government.



WWW SSL/TLS : Secure Socket Layer (HTTP + SSL == HTTPS) mail, data GPG/PGP : GNU privacy Guard / Pretty Good Privacy



The technical solution When there are no implementation errors, no incorrect security procedures (permissions, passwords), it is today possible to build a safe information transfer infrastructure over the Internet, for example for WWW applications (HTTP).

Unfortunately, it is still not possible to ensure the absence of implementation errors or incorrect security procedures, and this will not change in the foreseeable future. Not only the client workstations are subject to many and regular attacks (using security holes or incorrect default configuration for example in Microsoft software), but servers too are vulnerable : for example credit card numbers are usually stolen in mass through electronic merchants.



- encrypt/decrypt files (also symetric)
- sign files/verify signatures
- can be used together with e-mail software packages
- key servers, trust network

Main commands of GNU Privacy Guard (GPG)

generate a key pair	gpggen-key
generate a revocation certificate	gpggen-revoke
get key ID 7F76BFC9 from key server	gpgrecv-keys 7F76BFC9
send key ID 7F76BFC9 to key server	gpgsend-keys 7F76BFC9
get finger print from key	gpgfingerprint 7F76BFC9
sign a key with yours	gpgsign-key 7F76BFC9
verify a signature	gpgverify SIG-FILE FILE
encrypt	gpg -e < file > file.gpg
decrypt	gpg < file.gpg > file

Configuration of GNU Privacy Guard (GPG) The configuration file is

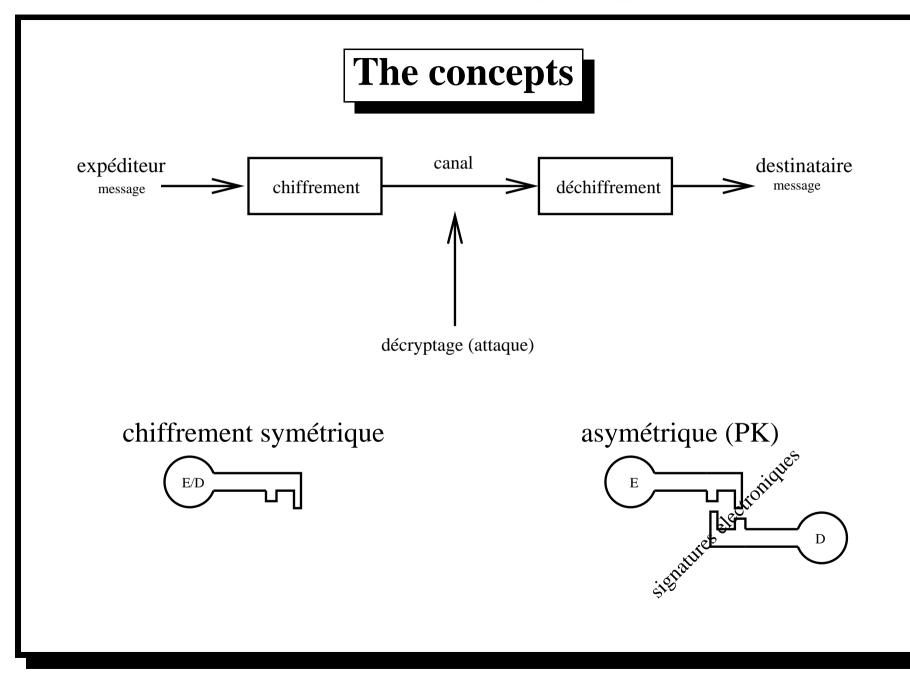
~/.gnupg/options

Some interesting options to set are :

load-extension idea Especially when handling older PGP keys. Warning : IDEA is patented ! Check the license of the package before using.

keyserver www.ch.pgp.net A key server to send/receive keys to/from. honor-http-proxy Set this if you have an HTTP proxy.

	Security Settings	
	PGP	
Identity	- Warning	
	Please check if encryption really works before you start using it seriously.	
Network	Also note that attachments are not encrypted by KMail.	
- CO	☐ Identity	
	PGP User Identity:	
Appearance	- Options -	
	Keep passphrase in memory	
Composer	Always encrypt to self	
()	X Show ciphered text after composing	
Mime Headers	Encryption tool	
\cap	Auto-detect	
	C GPG - Gnu Privacy Guard	
Security	C PGP Version 2.x	
103	C PGP Version 5.x	
Miscellaneous	C PGP Version 6.x	
a section of the		
Help	Default QK Apply C	ancel



Symetric cryptography (secret key) We use a key (secret, known to the emitter and the recipient of the message). This is infact a bijective transforming function which is to be applied to a message. The same function – or it's invert bijective function – is applied to decipher the message.

Properties :

- safe, well-known
- small key sizes are enough
- needs a pre-transmission (safe !) of the key
- key handling complex (one key per P2P link)

Asymetric cryptography (public key) Two keys are used (they are generated in a pair and are symetric/complement of each other). There is no *easy* way to go from one key to the other. One key can be used to cipher, and the other to decipher, and the other way. Only one of the keys is published, the other remains secret. Properties :

- slow
- no safe channel necessary to transmit keys
- makes electronic signatures possible.

makes key handling simpler (but without guaranteeing anything without a *trust network*).

In general, Asymetric cryptography is based on functions which are easy to compute in one direction and difficult in another. For example, take two small prime numbers, let's say 23 and 37. Multiply them together (that'a quite easy). You get 851. Now try to find them just from the product. A public key will be the 851. The private key will be the (23, 37). Properties on numbers will be used which require to know the factors to be able to decrypt.

In practice In practice, most of the used protocols combine symetric and asymetric technology. A *session key* is generated using the asymetric technology, and then this session key, transmitted safely through asymetric technology, is used to cipher the message, with a performance and safety advantage (difficult to break the asymetric keys since the data exchange is minimal, and the session key are not reused).



- Diffie-Hellman
- hashing
- signature
- random number generator
- key length (sym/asym)
- certificates

Diffie-Hellman This scheme, using properties of the modulo arithmetic, allows to exchange a key without previously establishing a safe channel, but without authentification.

Hashing and signatures A hash is a projection of a message to a shorter value. The value has differing properties depending on the hashing function. For example, changes in bits which are in sequence in the message will be detected until a certain number (this is usually used in the *Cyclical Redundancy Checks* in telecommunication). Or, slightly dissimilar messages generate a very different value and it's difficult to maintain a similar value artificially on a changed message (chaotical response). In a way, hashing functions are *trapdoor* functions : they cannot be inverted. The only way is to go through the space of all possible solutions. Because this space is usually big and hashing functions, when used for cryptographic purposes are designed chaotic, this is time consuming – and that's the goal. Examples :

– UNIX passwords

- md5sum UNIX command

Side effect : electronic signatures If a hash value is encrypted by the sender, using his *private key*, any recipient can decrypt it using the public key, and recompute the hash on the actual message. If the hashes are the same, the message was not modified. This works only, of course, if you know that the real sender's public key is the one you got. This question opens a can of worms : the *trust network* and certificates.

The problem of Trust

Problem : how do you know this is **my** public key?

Solution : *trust network* or Certificate Authorities



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Cnetennomatch%2Cnetenerror&fingerprint=on&exact=off&search=Jacques+Chirac 🗢
CRYPTO ABOUT ABOUT EMAIL THE CONTRACT ABOUT SERVER WEBMASTER
belgium.keyserver.net
 Primary Name or Identifier Secondary Name or Identifier Certificate (hold your mouse over a certificate to see its creation date)
SizeCreated L024 L999/0L/26
k472 3848 ildent.gouv.Er>
nated by <u>OKS</u> running at Belgium

from, as long as each of the signee trusts the signed key.

For example, with GPG or PGP, you can deposit your key(s) to a *key server*, for example http://www.keyserver.net along with the signatures of the person who directly know you. When you meet people and give them your key, they can sign it and deposit the signature on the key server directly.

It then suffices for any of the signee to be known directly or indirectly by a recipient and he will be able to assert that your key really comes from you.

Giving a finger-print of the key (a hash, really) is in general sufficient : by phone, or in person. In some cases you may want to check the ID card of the person.

All this is called a trust ring.

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cation et:11371/p	ks/lookup?op=vindex&template=netensearch%2Dnetennomatch%2Cnetenerror&fingerpr	int=on&exact=off&sea	rch=Debian+Switz
AD AK	EY CHOOSE ANOTHER CRYPTO ABOUT HES SERVER		
pekey IDName and Cer	tificates	SizeCreated	
	ebies Switzerland Distributor (debien-info@debien.ch) 5500 0093 715% 42%2 0960 P60% %900 002% 0868 6029	1024	2001/01/14
03086279	Marc SCHABPER (achaefer@alphanet.ch>		
7#76BPC9	Marc SCHARPER (gpg) KichaeferØalphanet.ch>		
08686029	Debian Switzerland Distributor <debian-info@debian.ch></debian-info@debian.ch>		
20851421	Marc SCHABPER Kachaefer@di.epfl.ch>		
29350089	Brik Rossen (Linux consultant) <rossen@freesurf.ch></rossen@freesurf.ch>		
A5080316	Frederic Schutz Kschutz@cui.unige.ch>		
28183815	Prederic Schutz Kachutz@mathgen.ch>		
РАСВОСЗС	Andreas Trottmann Kandreas.trottmann@werft22.com>		
9 14881838	Alexandre Galletet Kalex@galletet.ch>		
0068380	Ihsan Dogan Kgumboğnetlabs.org>		
22,923225	tresh.net Root Key (RSA) <pre>Kroot@tresh.net></pre>		
B0849#34	trash.net Root Key (055/DH) <root@trash.net></root@trash.net>		
7584P508	Thomas Bader <pre><pre>chomasb8trash.net></pre></pre>		
38487850	Thomas Bader <thomasbétrash.net></thomasbétrash.net>		
04950B1	Tobias Greach (Leletoby) <loby@fear.ch></loby@fear.ch>		
26168095	Istvan Sebestyen (GPG-KEY) Kstevee0alphanet.ch>		
42282PPB	<pre>PXOlix Hauri Informatician consultant) <felix@f-hauri.ch></felix@f-hauri.ch></pre>		

Certificate Authorities A certificate is a signed key. This signature means that the signee testifies that the key belongs to his owner. If you trust the signee, you trust the key !

All WWW browsers have a list of known-safe keys built-in. Unfortunately, providing a safe service is not a prerequisite from being listed : in general, you just need to give money to the WWW browser software editor.

This privilege doesn't mean that this authority doesn't deliver certificates to people who are not who they pretend.

For example, **Verisign** has delivered recently a certificate for the microsoft.com domain to a third-party (if you don't believe me, maybe you will believe Microsoft : http://support.microsoft.com/default.aspx?scid=kb;EN-US;q293818).

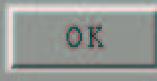
The real solution (until some reliable and trustworthy infrastructure exists, see above for Verisign) is to ask the finger-print of the key to the service provider directly, through a safer mean (for example telephone). Unfortunately, until recently at least,

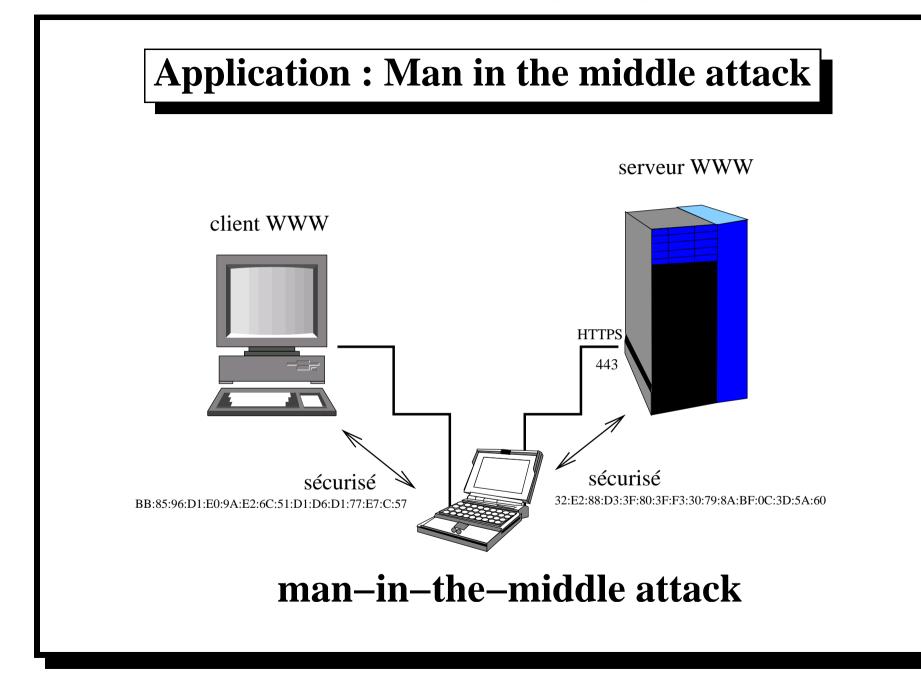
most of those didn't even know what a finger-print was.

() Netscape: View A Certificate X This Certificate belongs to: This Certificate was issued by: search.alphanet.ch search.alphanet.ch www.adm@alphanet.ch www.adm@alphanet.ch Search and DB services Search and DB services ALPHANET NF ALPHANET NF Colombier, Neuchatel, CH Colombier, Neuchatel, CH Serial Number: 00 This Certificate is valid from Sat Jul 14, 2001 to Tue Jul 09, 2002 **Certificate Fingerprint:** BB 85.96 D1 ED 9A E2 6C 51 D1 D6 D1 77 E7 9C 57



() Netscape: View A Certificate X This Certificate belongs to: This Certificate was issued by: www.verisign.com/CPS Incorp.by Ref. LIABILITY LTD.(c)97 www.yellowast.ch Pastfinance VenSign VeriSign International Server CA - Class 3 Die Schweizenische Post VenSign, Inc. Bern, Bern, CH VenSign Trust Network Serial Number: 34:20:0B:B0: A7:2E:94:71:E1:18: AA:E4:4B:E8:6A:9E This Certificate is valid from Thu Aug 16, 2001 to Sat Aug 17, 2002 **Certificate Fingerprint:** BC 19 89 9E 17 C4 F4 A9 F6 09 1B 9D 74 03 17 6D





Man in the middle attack This attack allows a cracker to read and modify all the data he wants in a connection to a secure WWW server. This of course would also work for e-mail in some cases. The attacker will simply reside on the path between the two computers and replace the key exchange by his own key, creating two safely encrypted tunnels instead of one. **DNS spoofing** attacks may make those attacks quite easy, even without a physical access to the paths between the two machines.

Solutions :

- trust rings
- certification authorities.

ask the key fingerprint directly to the service provider, through a safe channel.
In the case of WWW, note that in general, the WWW browsers will alarm if a server
SLL key changes : note that because keys usually have an expiration date (for security reasons), keys WILL change from time to time, making everything fuzzier.

Other attacks

by **decreasing importance** :

- social engineering
- software failures and bugs
- tempest radiations (wireless !)
- *brute force* (computers getting better)
- différential (DPA) (on-chip)
- theory attacks

weak link

Weak link As always, your problem is the weakest link : if a server offers SSL but it is not mandatory, or it ships passwords with e-mail in cleartext, it won't be safe.

The complexity of systems makes often human beings the weak link.



- ITU-T standard X.509
- specifies interoperable way to store and exchange certificates
- **OpenSSL** supports most operations



help build a stronger/wider trust network

Thank you !

More information

- Cours postgrade sécurité ESNIG, October 2002-June 2003, Neuchâtel, http://www.esnig.ch
- GPG Home Page : http://www.gnupg.org
- PKI Home Page: http://www.pki-page.org
- Crypto Home Page ETHZ : http://www.crypto.ethz.ch
- Thomas BADER's GPG article (German) :

http://www.linux-magazin.de/ausgabe/1999/12/GnuPG/gnupg.html

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